A proposal of new concurrency model for Ruby 3

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People love "Concurrency"





Improved scalability by relaxing the GVL Charlie Gracie @crgracie



ErRuby: Ruby on Erlang/OTP Lin Yu Hsiang @johnlinvc



concurrent-ruby and how it is making Rails concurrent Vipul A M @vipulnsward



EN

How to create multiprocess server on Windows with Ruby Ritta Narita @narittan



Ruby Concurrency compared Anil Wadghule @anildigital

EN



Why people love (to discuss) "Concurrency"?

- Performance by "Parallel" execution to utilize multiple-cores
- Ruby has thread system, but MRI doesn't permit to allow parallel execution.

About this presentation

- Show "Why difficult multi-threads programs"
- Propose new concurrent and parallel mechanism idea named "Guild"
 - For Ruby 3

Koichi Sasada

- A programmer living in Tokyo, Japan
 Ruby core committer since 2007
 YARV, Fiber, ... (Ruby 1.9)
 - •RGenGC, RincGC (Ruby 2...)



Koichi is an Employee

salesforce heroku

Difficulty of Multi-threads programming

Programming language evolution

- Trade-off: Performance v.s. Safety/Easily
 Performance: making faster programs
 - Safety: making bug-free programs
 - Easily: making programs with small efforts

Two example C language

- String manipulation with pointers
- Memory management without GC

String manipulation with pointers

- •C: Using raw pointers to manipulate strings
 - Good: all-purpose and fast
 - Bad: Error-prone
 - Generates strange behavior, such as abnormal termination
- Ruby: Wrap with String class
 - Good: Easy to use
 - Bad: slower than C in some cases

Object management without GC

- •C: Free memory objects manually
 - Good: full control (target, timing and so on)
 - Bad: Error-prone
 - double-free/memory-leak, ...
- Ruby: Automatic collection with GC
 - Good: nothing to care about object collection
 - Bad: introduce some overhead

Ruby chose "safety/easily" approach

- Ruby encourage "Happy Programming"
 - Reduce programmer's cost
 - Nowadays computer is enough faster
 - Implementation techniques overcome performance penalties

Do you want to program without GC?

Muilti-threads programming is difficult

- Introduce data race, race condition
- Introduce deadlock, livelock
 Difficult to make
- Difficulty on debugging because of correct (bug-free) nondeterministic behavior
 - difficult to reproduce same problem

• Difficult to tune performance

Difficult to make fast programs

Data race and race condition

- Bank amount transfer example
 - Quoted from Race Condition vs. Data Race http://blog.regehr.org/archives/490

def transfer1 (amount, account_from, account_to)
 if (account_from.balance < amount) return NOPE
 account_to.balance += amount
 account_from.balance -= amount
 return YEP
end</pre>

Data race

- "account_to.balance += amount" has Data-race
 - Assume two threads (T1 and T2) invoke this methods with same bank accounts

```
# interleave two threads (T1: amount = 100, T2: amount = 200)
T1: t1 = account_to.balance # t1 = 10
T2: t2 = account_to.balance # t2 = 10
T2: account_to.balance = t2 + 200 #=> 210
T1: account_to.balance = t1 + 100 #=> 110 (expected: 310)
```

Race condition

- To avoid data-race with the lock
- But there is another problem yet

Lock with "Thread.exclusive"
def transfer2 (amount, account_from, account_to)
if (account_from.balance < amount) return NOPE
<u>Thread.exclusive{</u> account_to.balance += amount }
<u>Thread.exclusive{</u> account_from.balance -= amount }
return YEP
end

Race condition

To avoid data-race with the lock
But there is another problem yet

- # T1 amount = 100, T2 amount = 200, account_from.balance = 250
- T1: if (account_from.balance (== 250) < 100) return NOPE # OK, go through
- T2: if (account_from.balance (== 250) < 200) return NOPE
- T2: Thread.exclusive{ account_to.balance += 200 }
- T2: Thread.exclusive{ account_from.balance -= 200 } #=> 250-200 => 50
- T1: Thread.exclusive{ account_to.balance += 100 }
- T1: Thread.exclusive{ account_from.balance -= 100 } #=> 50 100 => negative number!!

Final solution

Lock whole of method

def transfer1 (amount, account_from, account_to)
 <u>Thread.exclusive{</u>

if (account_from.balance < amount) return NOPE
 account_to.balance += amount
 account_from.balance -= amount
 return YEP
 }
end</pre>

Another example Multi-thread quiz

• What happen on this program?

```
ary = [1, 2, 3]
t1 = Thread.new{
    ary.concat [4, 5, 6]
}
t2 = Thread.new{
    p ary # what's happen?
}.join
```

```
(1) [1, 2, 3]
(2) [1, 2, 3, 4, 5, 6]
(3) (1) or (2)
```

Another example Multi-thread quiz

• Answer: (4) depends on an interpreter

ary = [1, 2, 3]
t1 = Thread.new{
 ary.concat [4, 5, 6]
}
t2 = Thread.new{
 p ary # what's happen?
}.join

On MRI, (3) is correct

It will shows [1, 2, 3] or [1, 2, 3, 4, 5, 6] (depends on thread switching timing)

Another example Multi-thread quiz

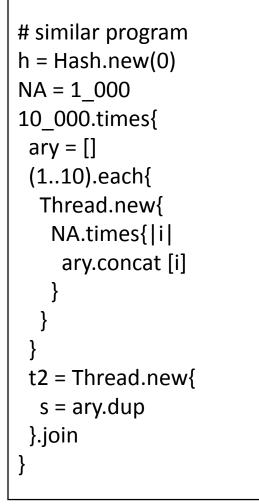
• Answer: (4) depends on an interpreter

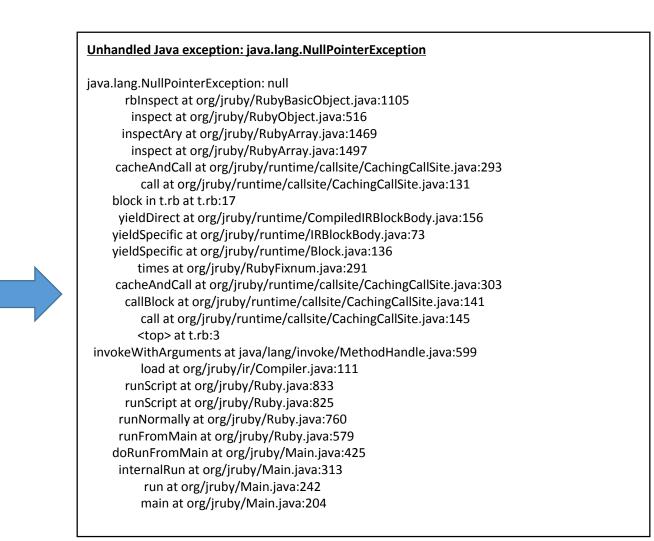
```
ary = [1, 2, 3]
t1 = Thread.new{
    ary.concat [4, 5, 6]
}
t2 = Thread.new{
    p ary # what's happen?
}.join
```

On JRuby:

It can cause Java exception because "Array#concat" is not thread safe

On JRuby ...





jruby 9.1.2.0 (2.3.0) 2016-05-26 7357c8f OpenJDK 64-Bit Server VM 24.95-b01 on 1.7.0_101-b00 +jit [linux-x86_64] On 8 hardware threads machine

Difficulty of multi-threads programs

- We need to synchronize all sharing mutable objects correctly
 - We need to know which methods are thread-safe.
 - Easy to track all on small program
 - Difficult to track on <u>big programs</u>, especially on programs using gems
- We need to check <u>all of source codes</u>, or believe <u>library documents</u> (but documents should be correct)
- Multi-threads prog. requires <u>"completeness"</u>

Difficulty of multi-threads programs (cont.)

- For debugging, it is difficult to find out the bugs
 - <u>Backtrace may not work</u> well because the problem may be placed on another line.
 - Bugs don't appear frequently with <u>small data</u>
 - Difficult to reproduce issues because of nondeterministic behavior

FYI:

Why MRI Array#concat is thread-safe?

- MRI uses GVL (Giant/Global VM Lock) to control thread switching timing and C methods (such as Array#concat) are working atomically.
- GVL prohibits parallel thread execution (BAD), however it avoids several severe issues (GOOD).

Thread programming: Performance tuning issue

a1 = []; a2 = [] $NA = 10 \ 000 \ 000$ t1 = Thread.new{ NA.times{|i| a1 << i } }.join t2 = Thread.new{ NA.times{|i| a2 << i} }.join

Serial program:

real 0m8.568s

user 0m37.816s sys 0m5.530s on JRuby

```
Thread programming:
Performance tuning issue
```

a1 = []; a2 = [] NA = 10_000_000 t1 = Thread.new{ NA.times{|i| a1 << i } } t2 = Thread.new{

NA.times{|i| a2 << i }

t1.join; t2.join

Parallel program

(2 threads):

real 0m6.411s

user 0m20.527s

sys 0m7.798s

Thread programming: Performance tuning issue

a1 = []; a2 = []

NA = 10_000_000

```
m1, m2 = Mutex.new, Mutex.new
```

t1 = Thread.new{

NA.times{|i| m1.synchronize{ a1 << i }}

t2 = Thread.new{ NA.times{|i| <u>m2.synchronize</u>{ a2 << i }}

t1.join; t2.join

Parallel program with a useless lock 1

(2 threads):

real0m10.264suser0m38.370ssys0m4.406s

Thread programming: Performance tuning issue

a1 = []; a2 = []

```
NA = 10_000_000
```

```
m = Mutex.new
```

```
t1 = Thread.new{
```

```
NA.times{|i| <u>m.synchronize</u>{ a1 << i }}
```

t2 = Thread.new{ NA.times{|i| <u>m.synchronize</u>{ a2 << i }}

t1.join; t2.join

Parallel program with a useless lock 2

(2 threads):

real0m15.163suser0m45.317ssys0m9.658s

Performance tuning issue

	Execution time
Serial program	<u>8.568s</u>
Parallel program	<u>6.411s</u>
Parallel program with a useless lock 1	<u>10.264s</u>
Parallel program with a useless lock 2	<u>15.163s</u>

Thread programming: Performance tuning issue

We need to use just correct number locks

Not enough \rightarrow unexpected behaviorToo much \rightarrow performance penalty

FYI: synchronization mechanism

- Many synchronization mechanisms...
 - Mutual exclusion (Mutex), monitor, critical section
 - Transactional memory (optimistic lock)
 - Atomic instructions
 - Synchronized Queue
 - ...
 - Research on many lightweight lock algorithms
- They assume we can use them correctly

Overcome thread difficulty



Problem:

Easy to share mutable objects

Idea:

Do not allow to share mutable objects without any restriction

Study from other languages

- Shell script with pipes, Racket (Place)
 - Copy mutable data between processes w/ pipes
- Erlang/Elixir
 - Do not allow mutable data
- Clojure
 - Basically do not allow mutable data
 - Special data structure to share mutable objects
 - Note that it can share mutable objects on Java layer NOTE: we do not list approaches using "type system"

Don't you know Elixir language?

Programming Elixir 1.2 by Dave Thomas

邦訳: プログラミングElixir 笹田耕一・鳥井雪共訳 2016/08/19

You can buy it TODAY!! サイン会は明日13時らしいです





Summary of approaches

- Communication with copied data (shell scripts)
 - Good: we don't need locks
 - Bad: copy everything is <u>slow</u>
- Prohibit mutable objects
 - Good: we don't need locks
 - Bad: Ruby utilizes many "write" operations. Unacceptable.
- Provide special data structure to share mutable objects
 - Good: we don't need locks (who don't use such special data structures)
 - Bad: Difficult to use special data structures.

Background was finished

Our goal for Ruby 3

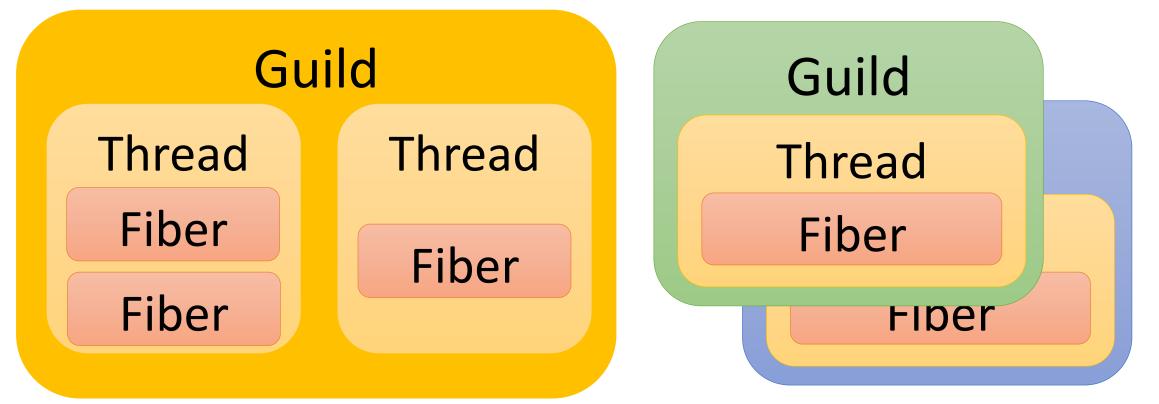
- We need to keep compatibility with Ruby 2.
- We can make **parallel program**.
- We shouldn't consider about locks any more.
- We <u>can share</u> objects with copy, but <u>copy</u> <u>operation should be fast.</u>
- We should share objects if we can.
- We can **provide special objects** to share mutable objects like Clojure if we really need speed.

"Guild"

New concurrency model for Ruby 3

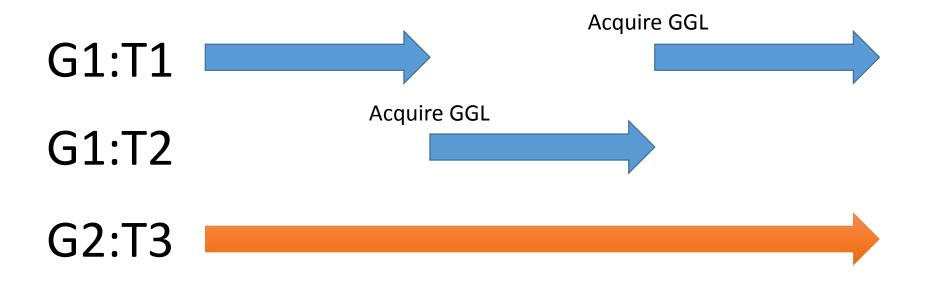
Guild: New concurrency abstraction

 Guild has at least one thread (and a thread has at least one fiber)



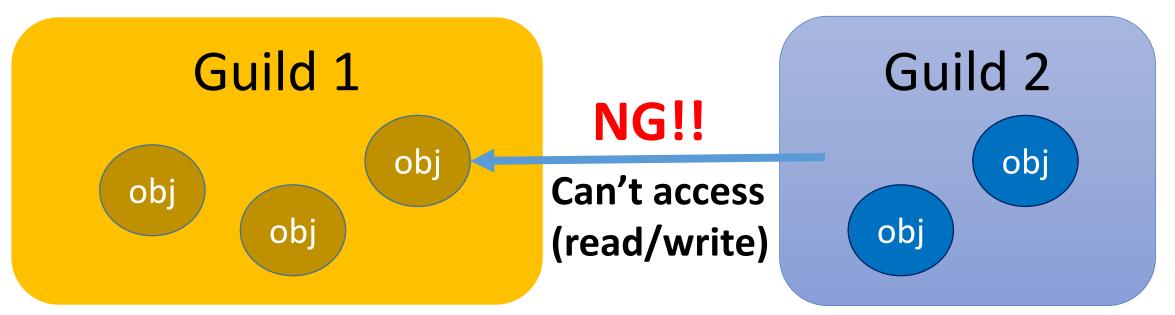
Threads in different guilds can run in Parallel

- Threads in different guilds <u>can run in parallel</u>
- Threads in a same guild <u>can not run in parallel</u> because of GVL (or GGL: Giant Guild Lock)



Guild and objects:

- All objects have their own membership
- All of mutable objects should belong to only one Guild (all mutable objects are member of one guild)
- Other guilds can not access objects



Object membership

Only one guild can access mutable object
→ We don't need to consider about locks

Because: NO data races and NO race conditions (if all guilds use only one thread)

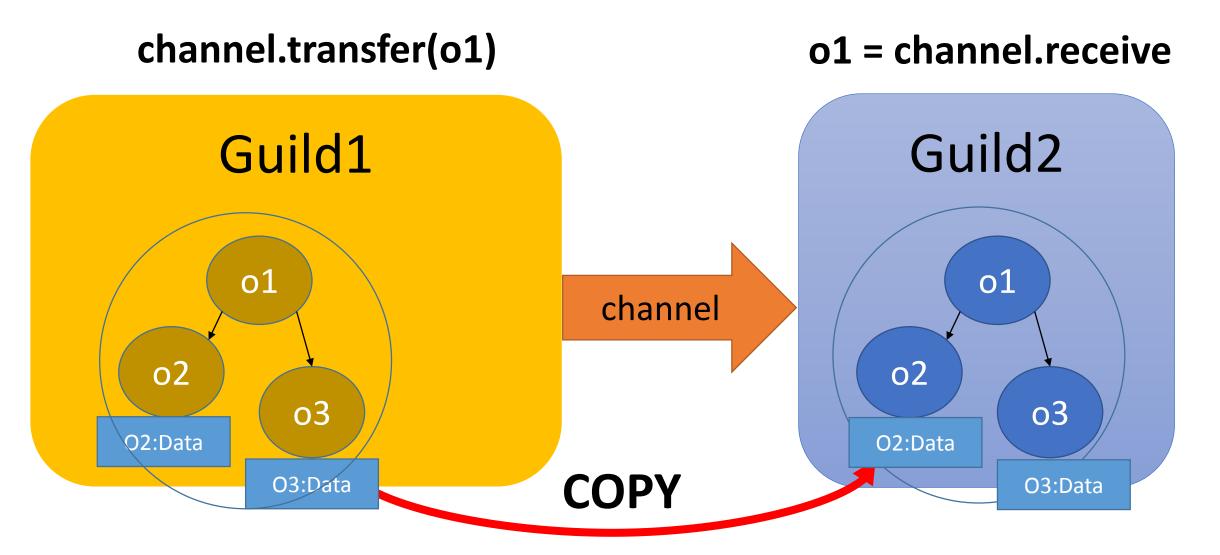
Inter guilds communication

- •"Guild::Channel" to communicate each guilds
- Two communication methods
 - **1.** Copy
 - 2. Transfer membership or Move in short

Copy using Channel

- Guild::Channel#transfer(obj) send <u>deep copied</u> object(s) to a destination guild.
- dRuby and multi-process system use this kind of communication

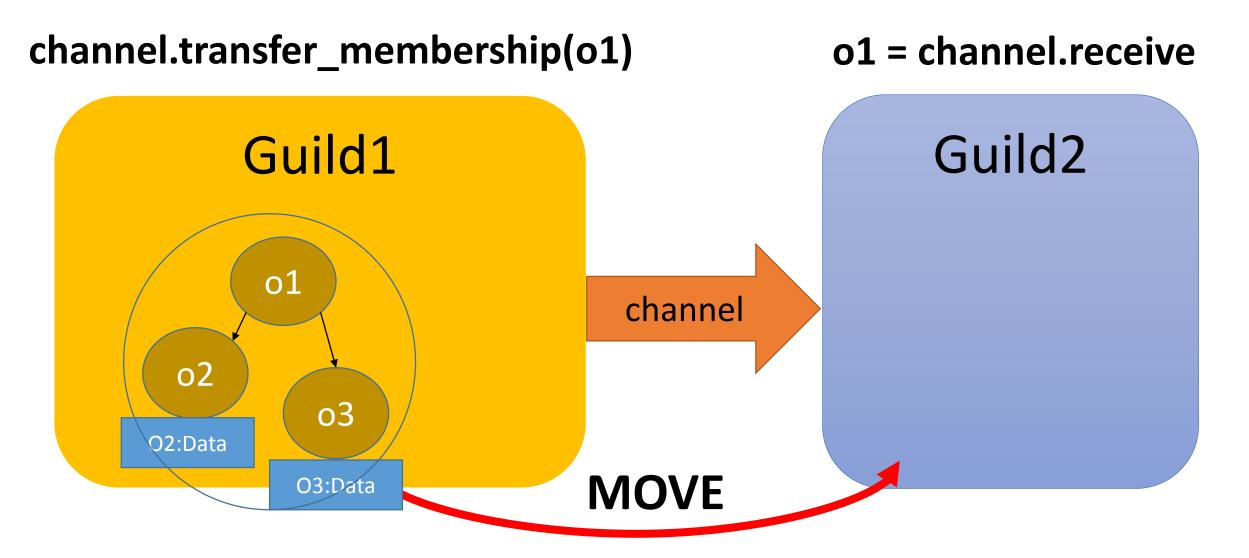
Copy using Channel



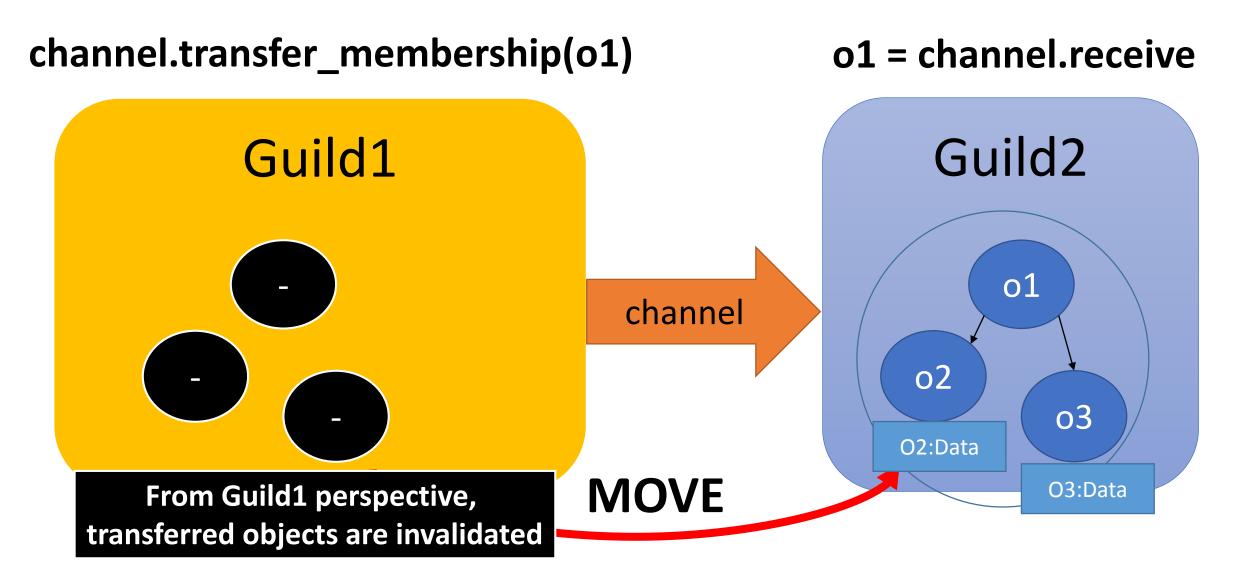
Move using Channel [New technique!!]

- Guild::Channel#transfer_membership(obj) change the membership of object(s)
 - Leave from the source guild
 - Join to the destination guild
- Prohibit accessing to left objects
 - Cause exceptions and so on

Move using Channel

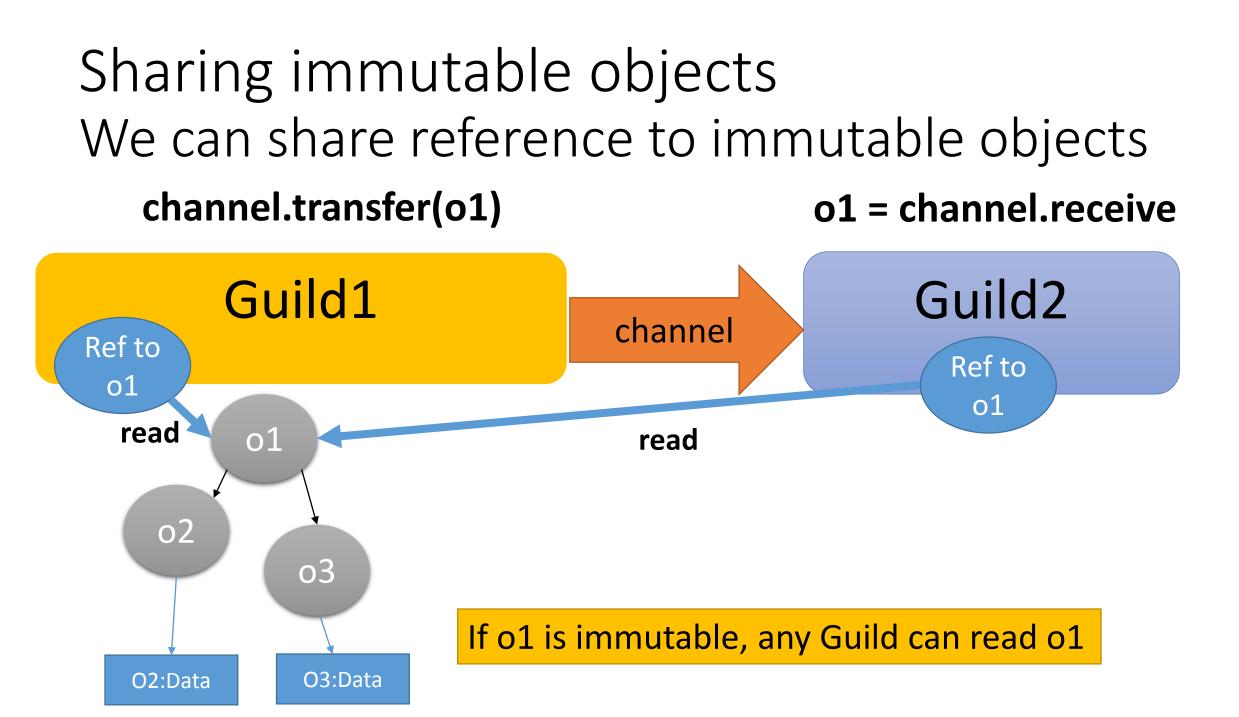


Move using Channel



Sharing immutable objects

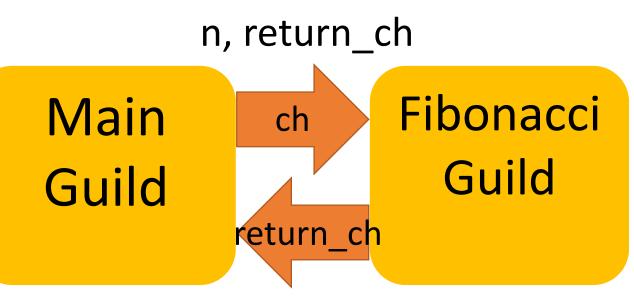
- Immutable objects can be shared with any guilds
 - a1 = [1, 2, 3].freeze: a1 is **Immutable object**
 - a2 = [1, Object.new, 3].freeze: a2 is not immutable
- We only need to send references
 very lightweight, like thread-programming
- •Numeric objects, symbols, true, false, nil are immutable (from Ruby 2.0, 2.1, 2.2)



Use-case 1: master – worker type

def fib(n) ... end
g_fib = Guild.new(script: %q{
 ch = Guild.default_channel
 while n, return_ch = ch.receive
 return_ch.transfer fib(n)
 end
})

ch = Guild::Channel.new
g_fib.transfer([3, ch])
p ch.receive

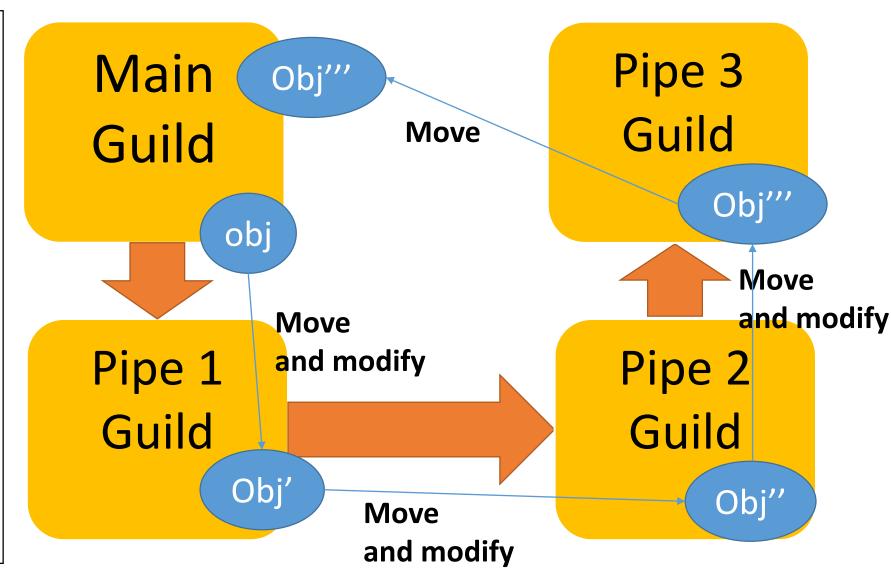


Answer of fib(n)

NOTE: Making other Fibonacci guilds, you can compute fib(n) in parallel

Use-case 2: pipeline

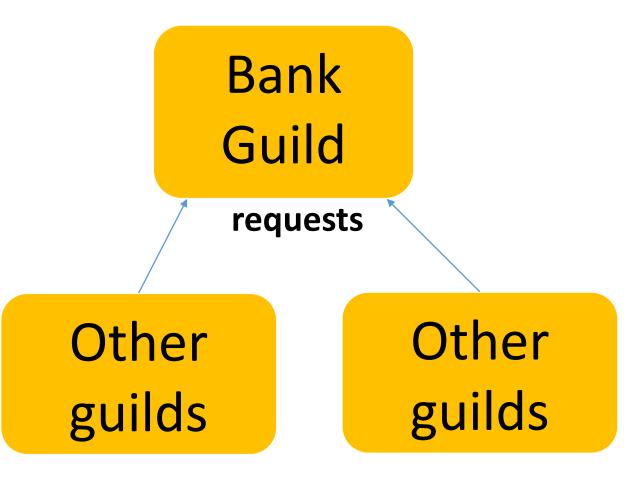
result ch = Guild::Channel.new g pipe3 = Guild.new(script: %q{ while obj = Guild.default channel.receive obj = modify obj3(obj) Guild.argv[0].transfer membership(obj) end }, argv: [result ch]) g pipe2 = Guild.new(script: %q{ while obj = Guild.default channel.receive obj = modify_obj2(obj) Guild.argv[0].transfer membership(obj) end }, argv: [g pipe3]) g_pipe1 = Guild.new(script: %q{ while obj = Guild.default_channel.receive obj = modify obj1(obj) Guild.argv[0].transfer_membership(obj) end }, argv: [g_pipe2]) obj = SomeClass.new g pipe1.transfer membership(obj) obj = result ch.receive



Use-case: Bank example

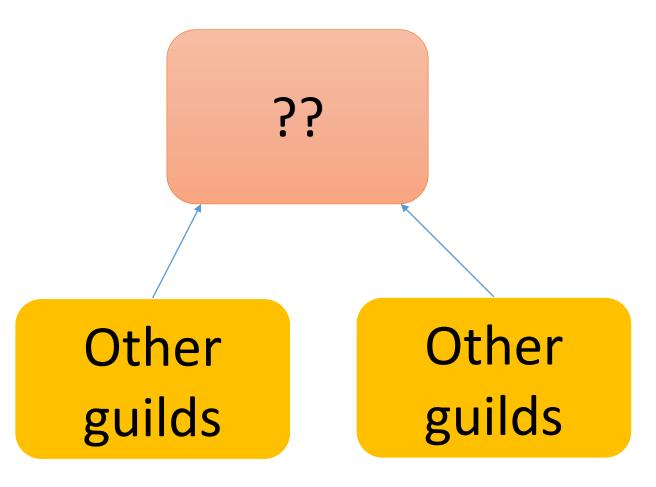
```
g_bank = Guild.new(script: %q{
 while account_from, account_to, amount,
       ch = Guild.default_channel.receive
  if (Bank[account_from].balance < amount)</pre>
   ch.transfer :NOPE
  else
   Bank[account_to].balance += amount
   Bank[account_from].balance -= amount
   ch.transfer :YEP
  end
 end
...
```

Only bank guild maintains bank data



Use-case: Introduce special data structure

- Ideas of special data structure to share mutable objects
 - Use external RDB
 - In process/external Key/value store
 - Software transactional memory



Summary of use cases

- Making multiple workers and compute in parallel
 - Requests and responses are communicate via channels
 - You can send it with copy or move
 - Maybe web application can employ this model
- Making Pipeline structures and compute in parallel
 - Each task has own Guild
 - Receive target object, modify it and send it next pipeline
 - You will send it with move (transfer membership)
 - It will help applications like applying several filters for input data
- Own responsibility by one Guild
 - All accesses are managed by one responsible Guild
 - If you want to share mutable objects, we need special data structures
 - External RDBs or key/value stores are also good idea for this purpose

Communication strategy

[Upper is better]

- Passing immutable objects
- Copy mutable objects
- If you have performance problem, move (transfer membership) mutable objects
- If you have performance problem too, use special data structure to share mutable objects

Compare between

- Thread model and Guild model
- •On threads, it is <u>difficult to find out</u> which objects are shared mutable objects
- On Guilds, there are no shared mutable objects
 - If there are special data structure to share mutable objects, we only need to check around this code

→ Encourage "Safe" and "Easy" programming

Compare between

- Thread model and Guild model
- On threads, inter threads communication is very fast.
- On guilds, inter guilds communication introduce overhead
 - "Move" (transfer membership) technique can reduce this kind of overheads

Trade-off: Performance v.s. Safety/Easily Which do you want to choose?

Digression: The name of "Guild"

- "Guild" is good metaphor for "object's membership"
- Check duplication
 - First letter is not same as other similar abstractions
 - For variable names
 - P is for Processes, T is for Threads, F is for Fibers
 - There are no duplicating top-level classes and modules in all of rubygems

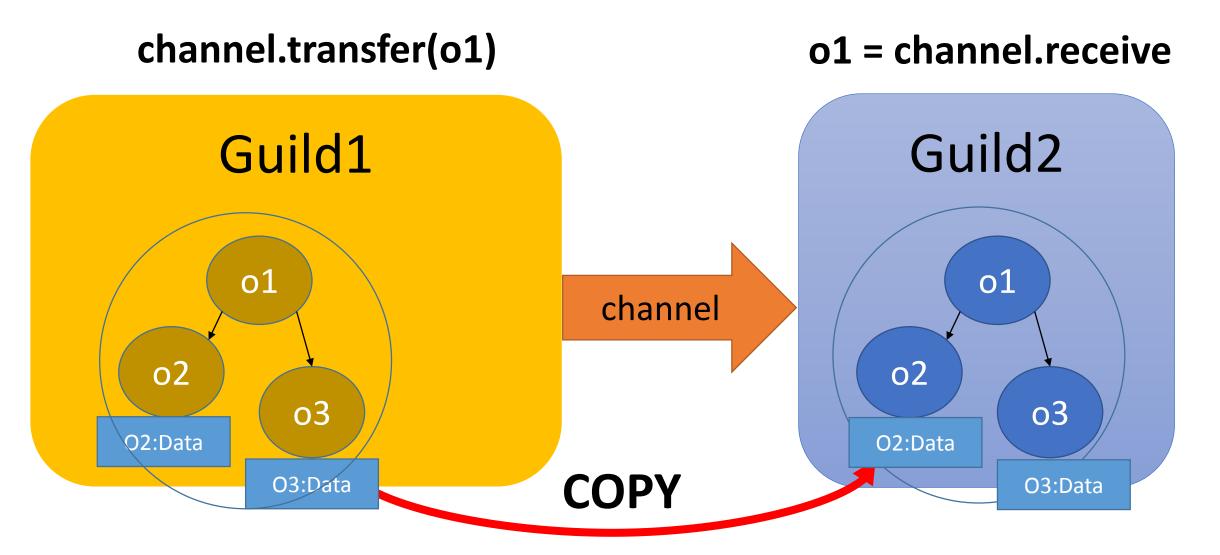
Implementation of "Guild"

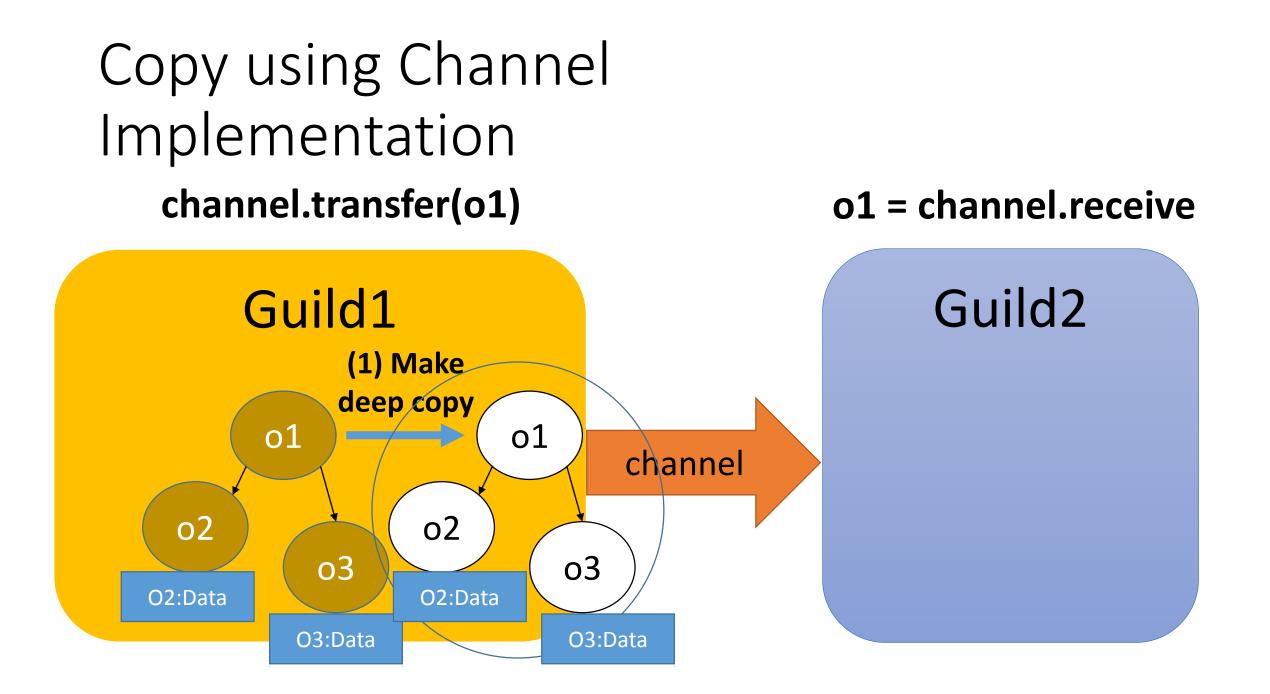
- How to implement inter Guilds communication
- How to isolate process global data

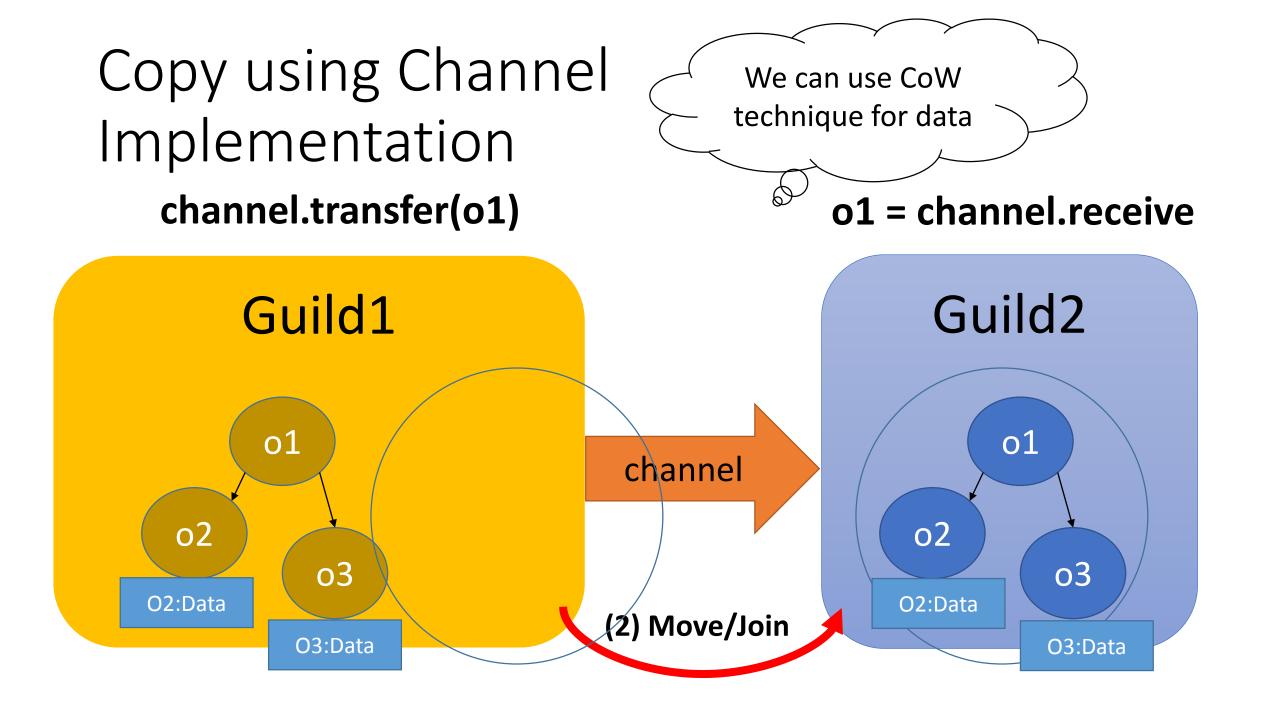
How to implement inter Guilds communication

- Copy
- Move (transfer membership)

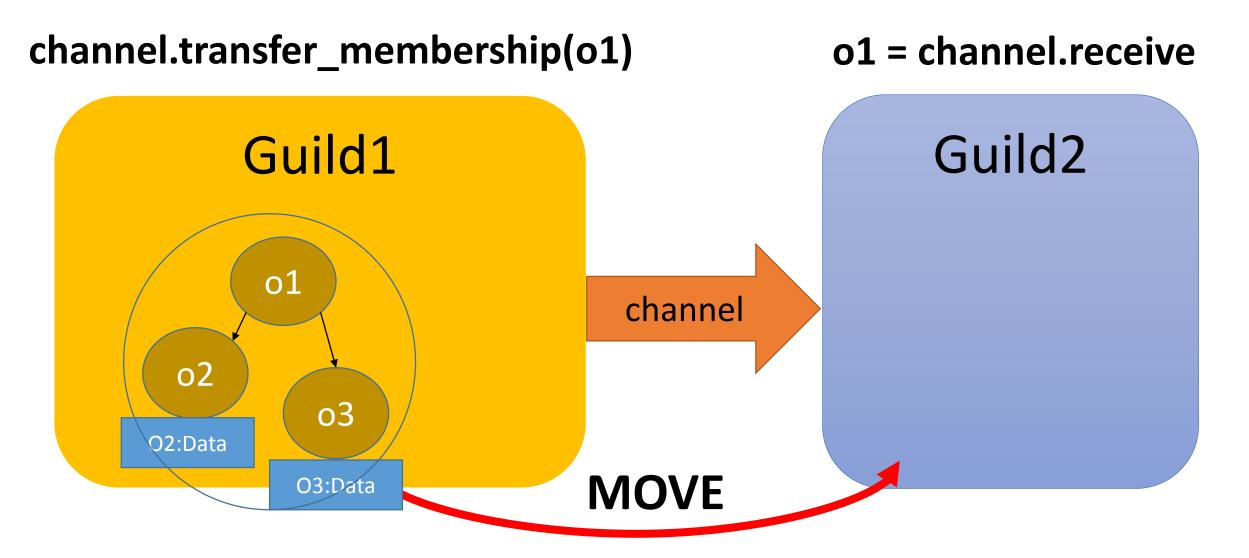
Copy using Channel



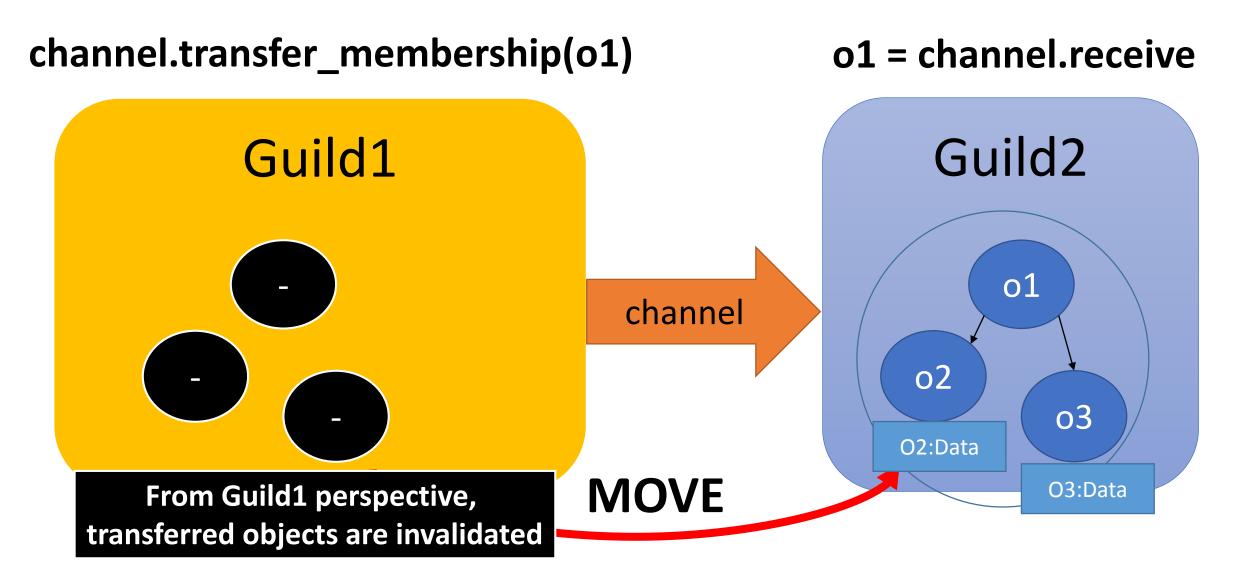


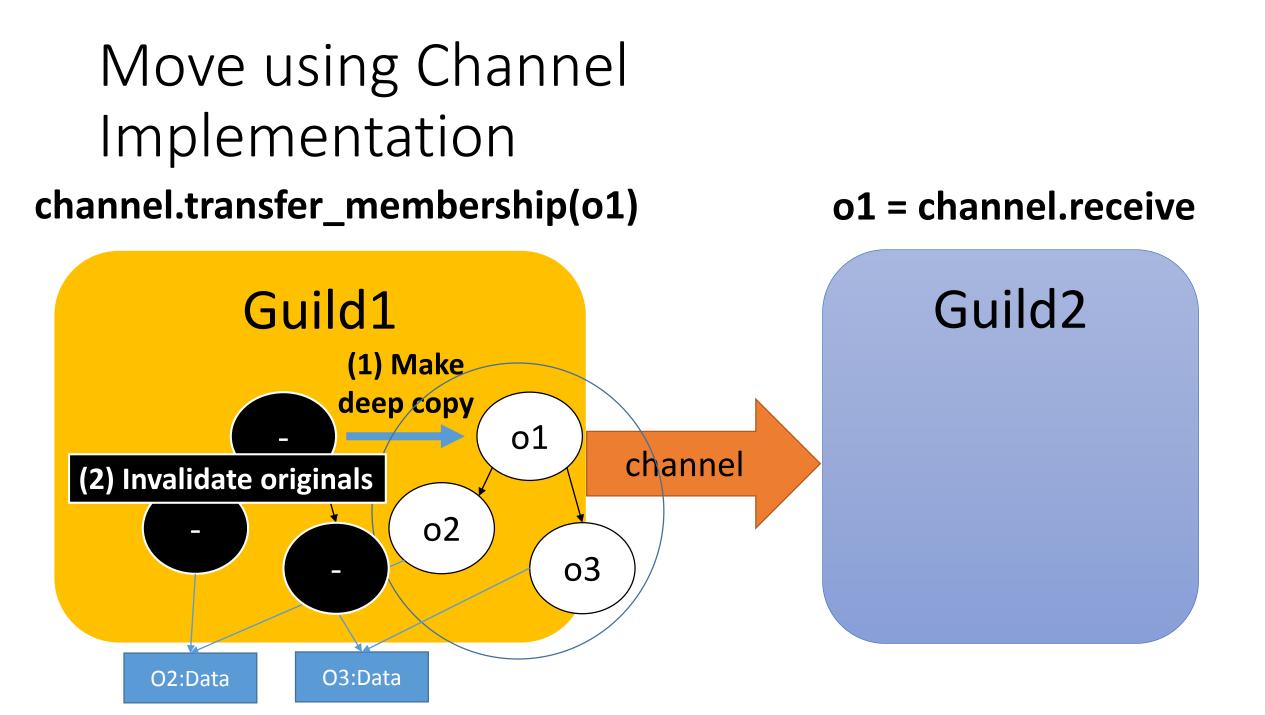


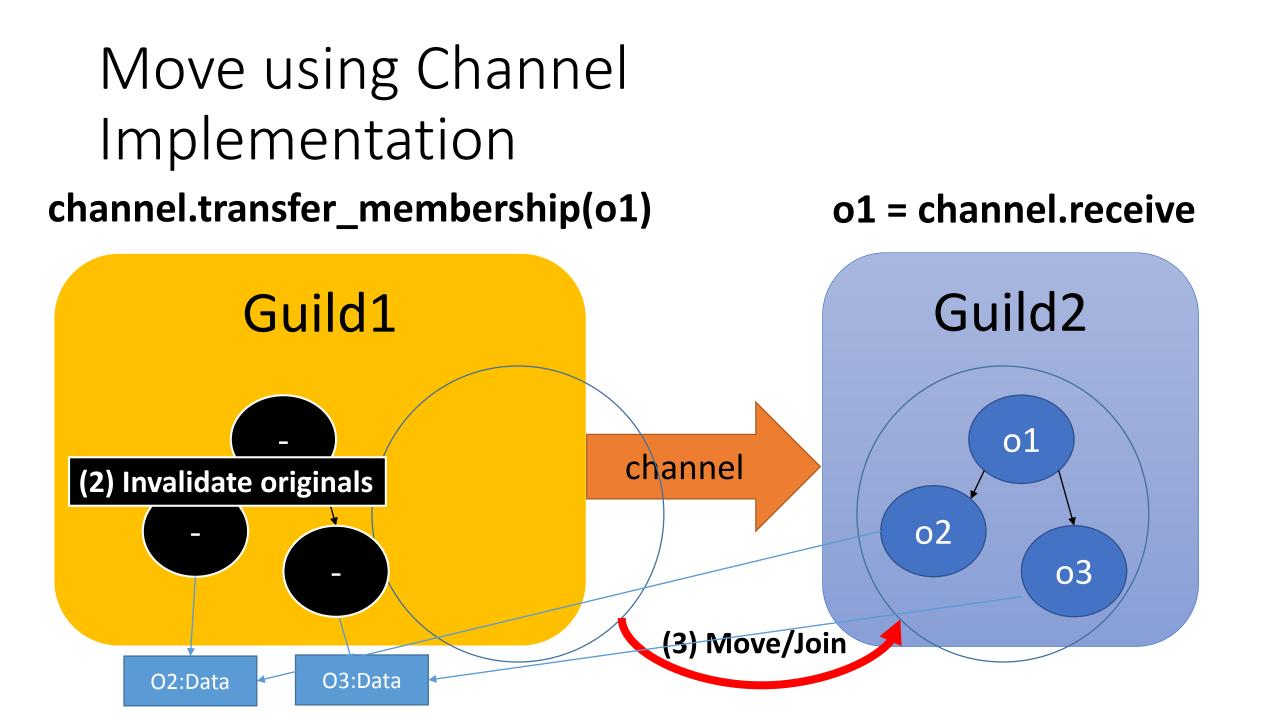
Move using Channel



Move using Channel







Ruby global data

- Global variables (\$foo)
 - Change them to Guild local variables
- Class and module objects
 - Share between guilds
- Class variables
- Keep compatibility with Ruby 2 • Change them to guild local. So that it is guild/class local variables
- Constants
 - Share between guilds
 - However if assigned object is not a immutable object, this constant is accessed only by setting guilds. If other guilds try to access it, them cause error.
- Instance variables of class and module objects
 - Difficult. There are several approaches.
- Proc/Binding objects
 - Make it copy-able with env objects or env independent objects
- ObjectSpace.each object
 - OMG

Interpreter process global data

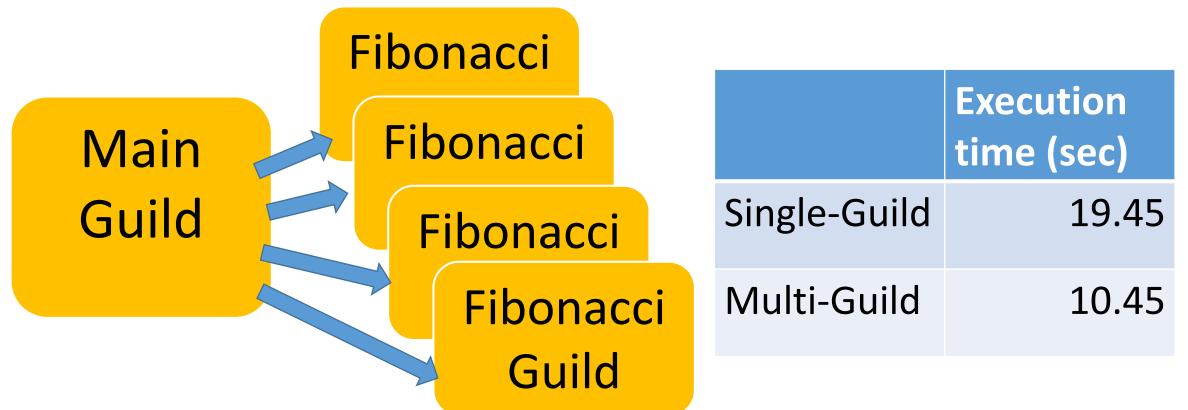
• GC/Heap

- Share it. Do stop the world parallel marking- and lazy concurrent sweeping.
- Synchronize only at page acquire timing. No any synchronization at creation time.
- Inline method cache
 - To fill new entry, create an inline cache object and update atomically.
- Tables (such as method tables and constant tables)
 - Introduce mutual exclusions.
- Current working directory (cwd)
 - Each guild should have own cwd (using openat and so on).
- Signal
 - Design new signal delivery protocol and mechanism
- C level global variables
 - Avoid them.
 - Main guild can use C extensions depends on them
- Current thread
 - Use TLS (temporary), but we will change all of C APIs to receive context data as first parameter in the future.

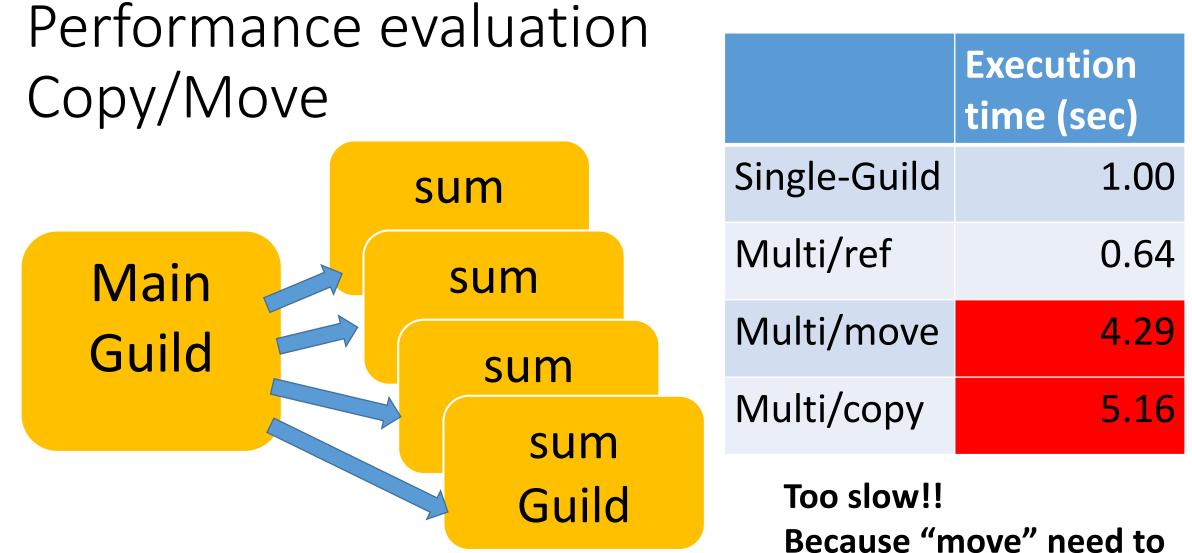
Performance evaluation

- On 2 core virtual machine
 - Linux on VirtualBox on Windows 7
- Now, we can't run Ruby program on other than main guild, so other guilds are implemented by C code

Performance evaluation Simple numeric task in parallel

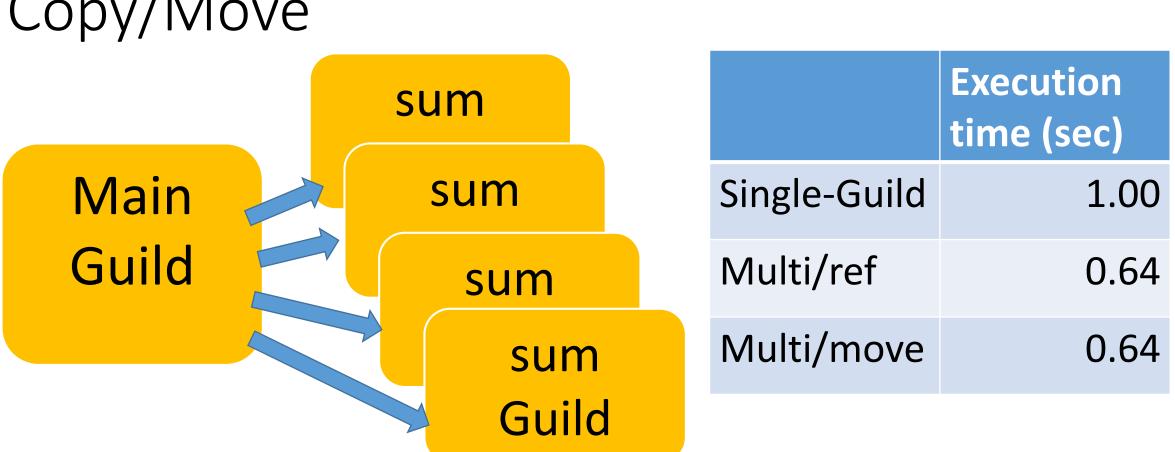


Total 50 requests to compute fib(40) Send 40 (integer) in each request



Total 100 requests to compute sum of array Send (1..10_000_000).to_a in each request

Because "move" need to check all of elements



If we know this array only has immutable objects, we don't need to check all elements => special data structure

Performance evaluation Copy/Move

Check our goal for Ruby 3

- Satisfied • We need to keep compatibility with Ruby 2.
 - OK: Only in main guild, it is compatible.
- We can make **parallel program**.
 - OK: Guilds can run in parallel.
- We **shouldn't consider** about locks any more.
 - OK: Only using copy and move, we don't need to care locks.
- We can share objects with copy, but copy operation should be fast.
 - **OK:** Move (transfer membership) idea can reduce overhead.
- We **should share objects** if we can.
 - **OK:** We can share immutable objects fast and easily.
- We can provide special objects to share mutable objects like Clojure if we really need speed.
 - OK: Yes, we can provide.

Summary

- Introduce "why threads are very difficult"
- Propose new concurrency abstraction "Guild" for Ruby 3
 - Not implemented everything yet, but I show key ideas and preliminary evaluation

Thank you for your attention

Koichi Sasada <ko1@heroku.com>

Approach comparison

	Process/MVM	Place (Racket)	Guild (copy/move)	Thread
Неар	Separate	Separate	Share	Share
Communication Mutable objects	Сору	Сору	Copy/Move	Share
Communication Frozen object	Сору	Share (maybe)	Share	Share
Lock	Don't need	Don't need	(mostly) Don't need	Required
ISeq	Сору	Share	Share	Share
Class/Module	Сору	Copy (fork)	Share	Share